



# Chapter 19

*(continued)*

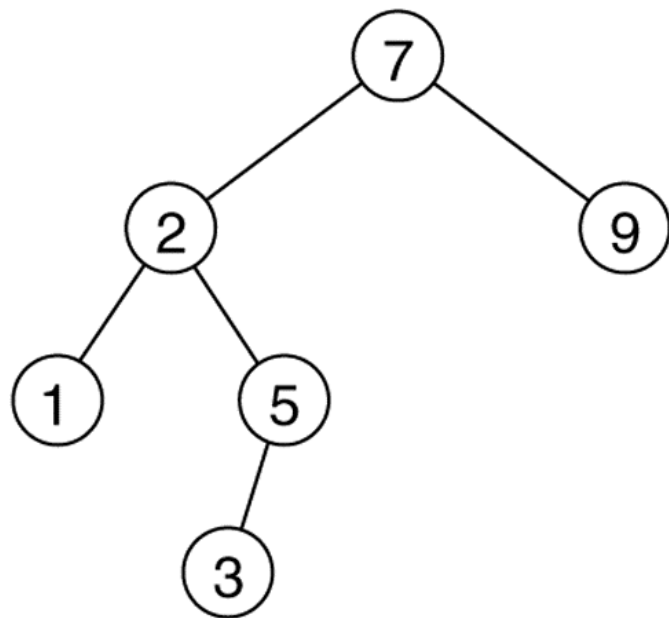
## Binary Search Trees

Data Structures & Problem Solving  
Using JAVA  
Second Edition

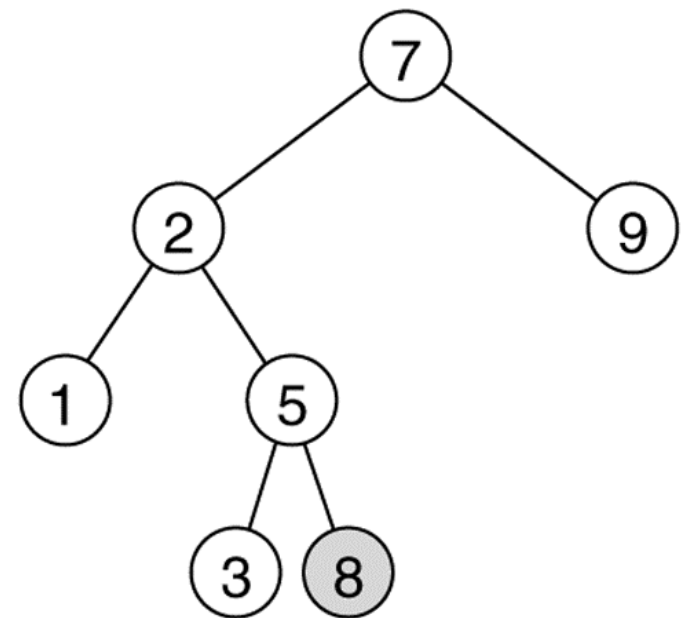
Mark Allen Weiss

# Figure 19.1

Two binary trees: (a) a search tree; (b) not a search tree



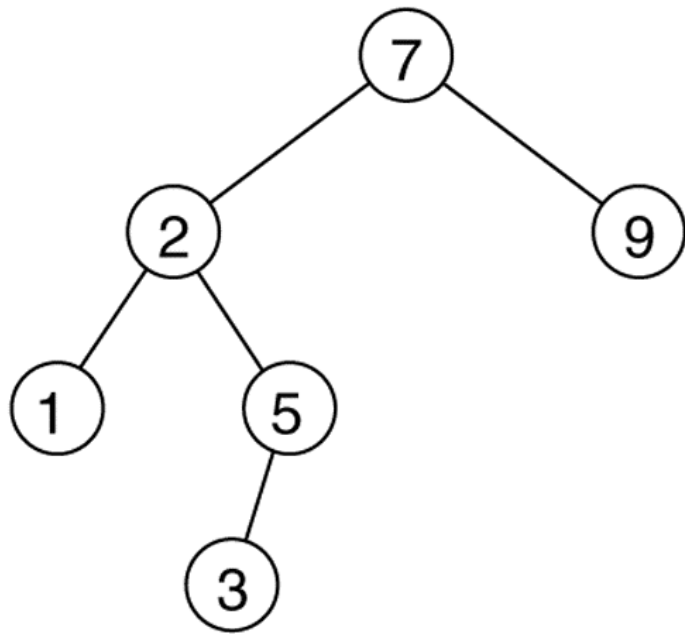
(a)



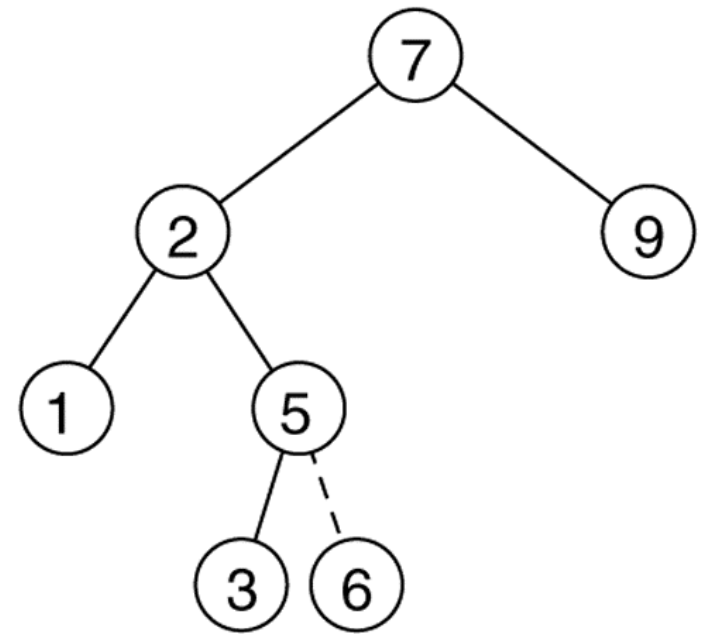
(b)

# Figure 19.2

Binary search trees (a) before and (b) after the insertion of 6



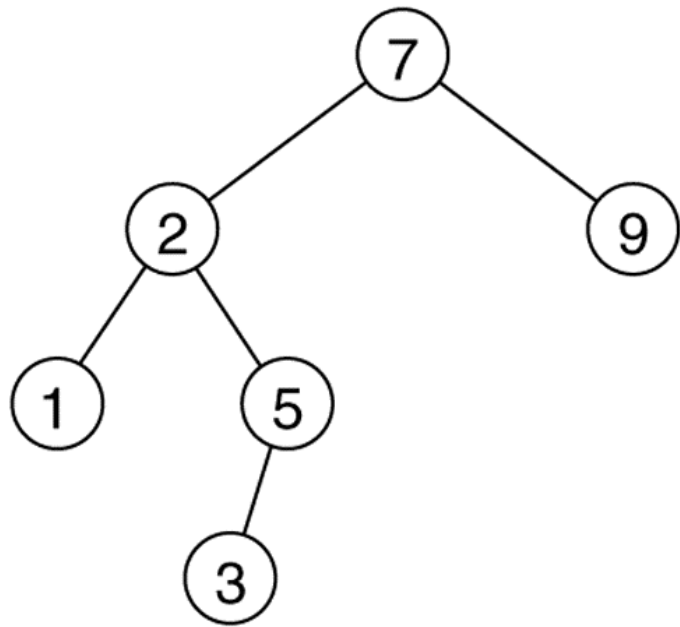
(a)



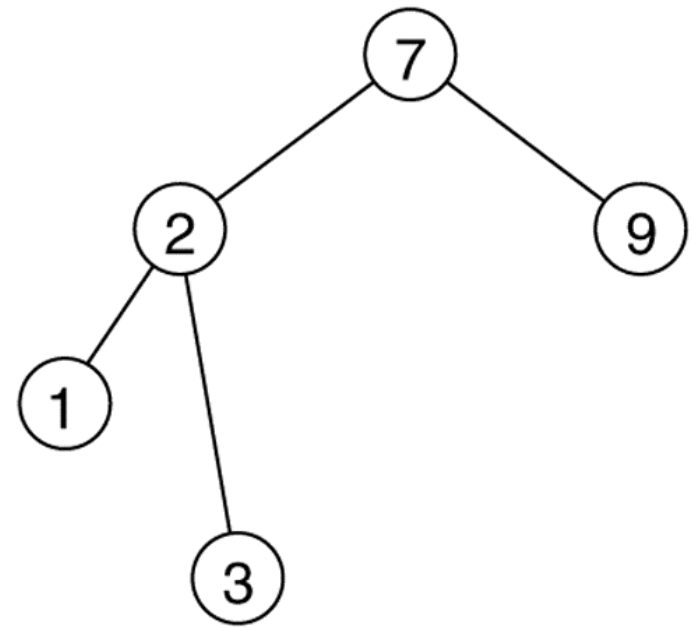
(b)

## Figure 19.3

Deletion of node 5 with one child: (a) before and (b) after



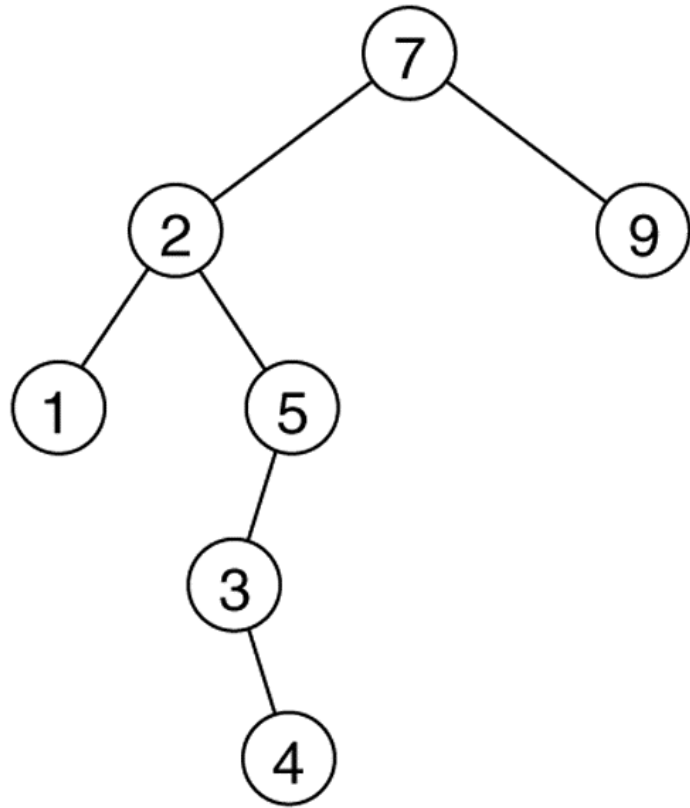
(a)



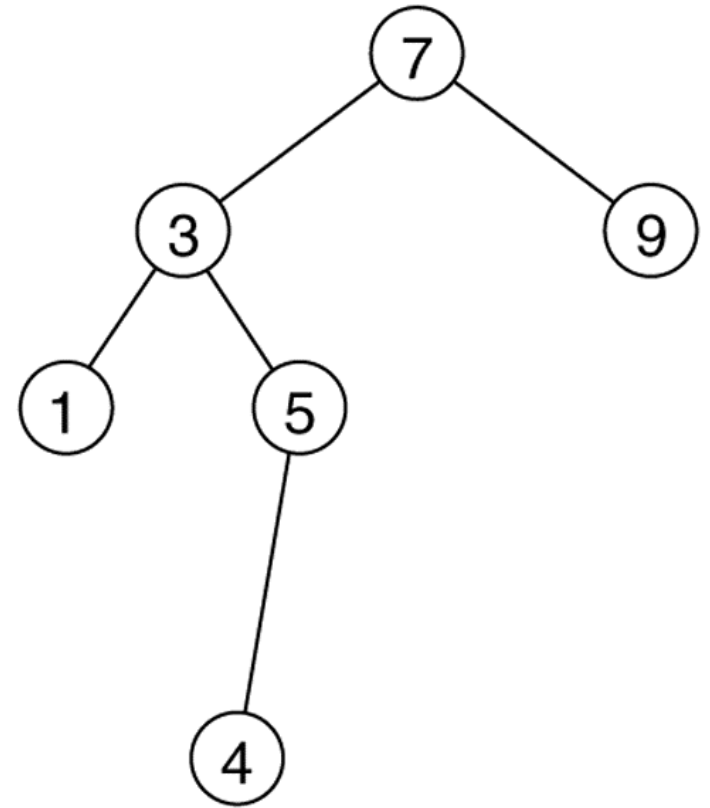
(b)

# Figure 19.4

Deletion of node 2 with two children: (a) before and (b) after



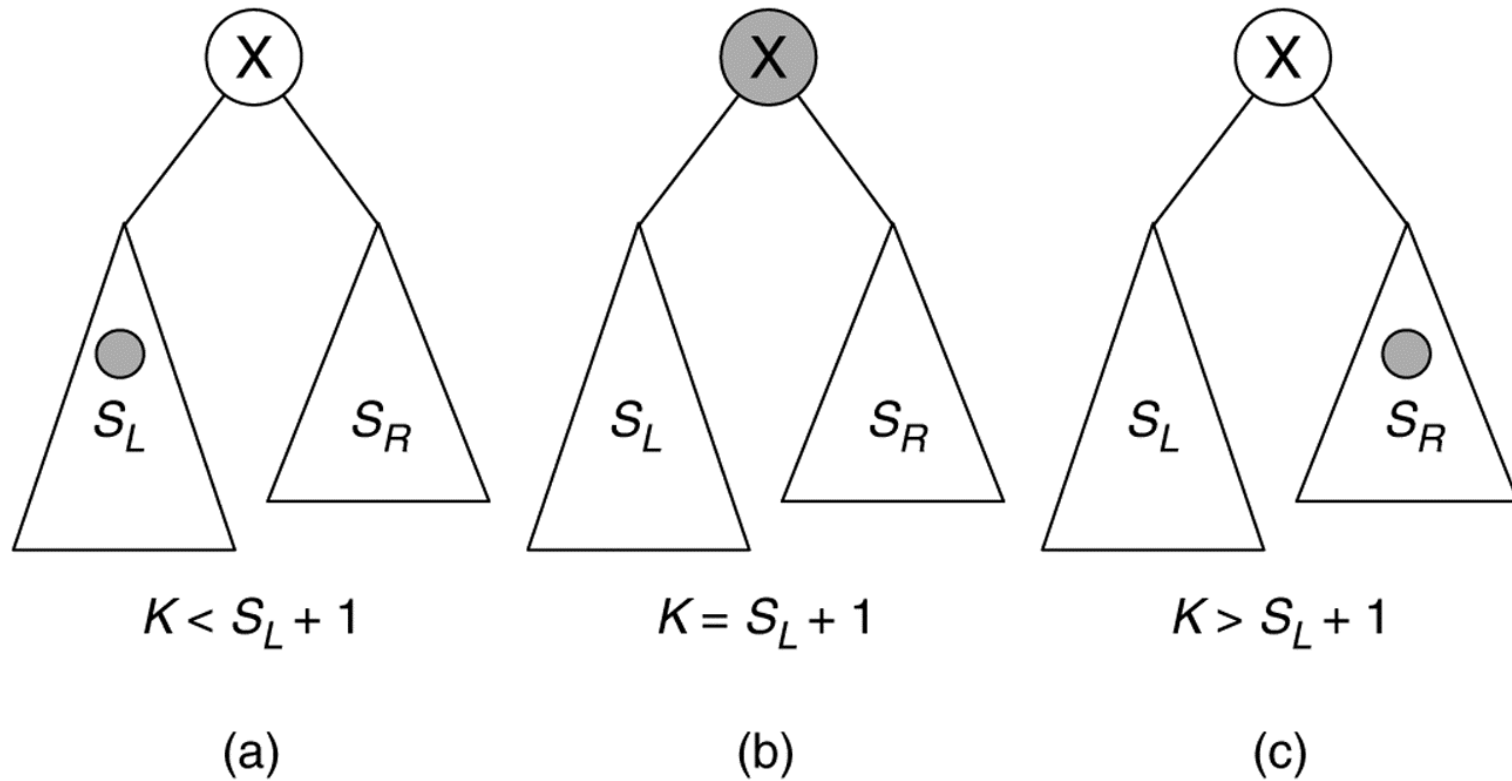
(a)



(b)

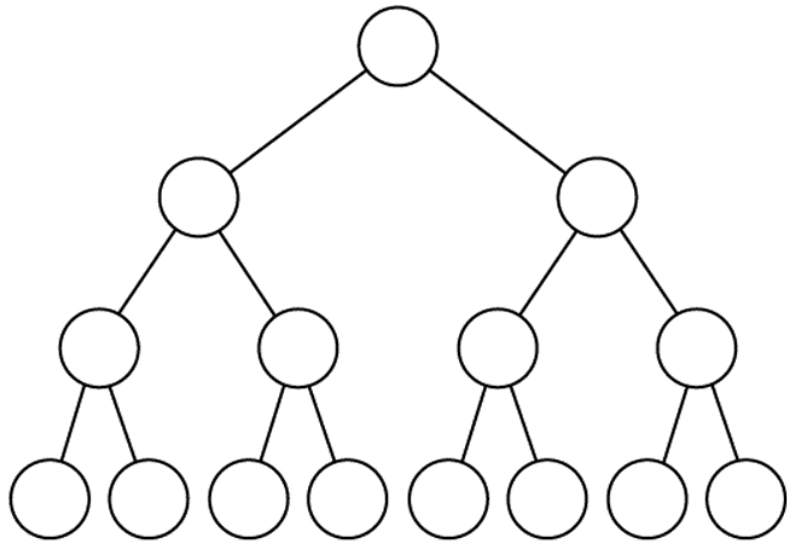
# Figure 19.13

Using the size data member to implement findKth

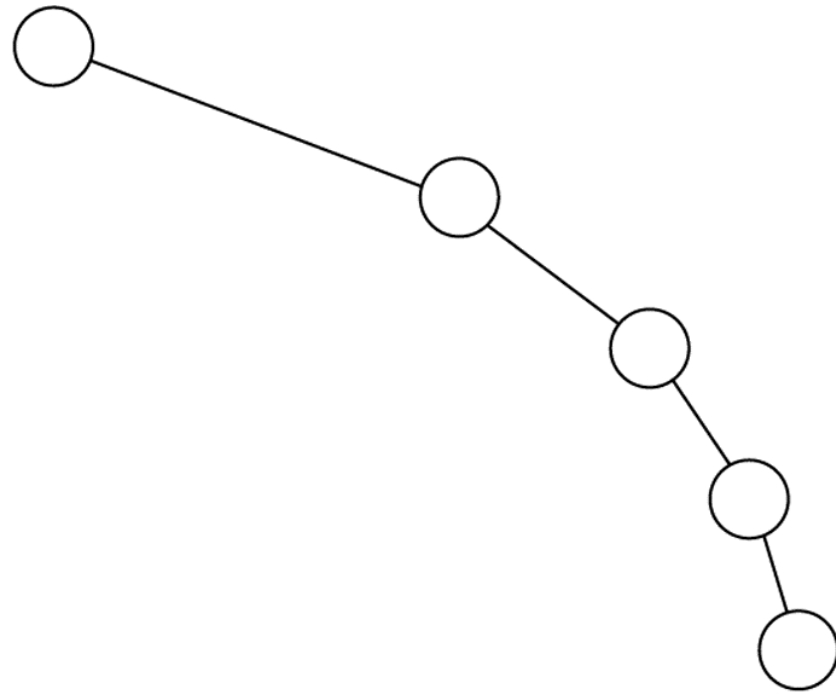


## Figure 19.19

(a) The balanced tree has a depth of  $\log N$ ; (b) the unbalanced tree has a depth of  $N - 1$ .



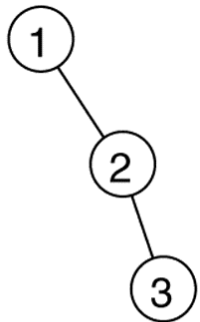
(a)



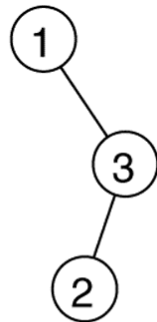
(b)

## Figure 19.20

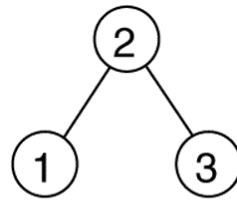
Binary search trees that can result from inserting a permutation 1, 2, and 3; the balanced tree shown in part (c) is twice as likely to result as any of the others.



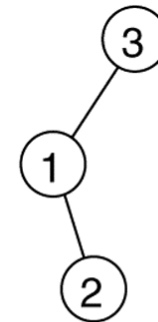
(a)



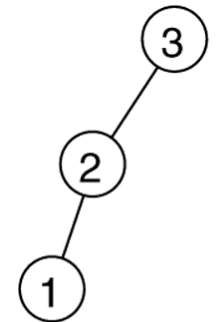
(b)



(c)



(d)

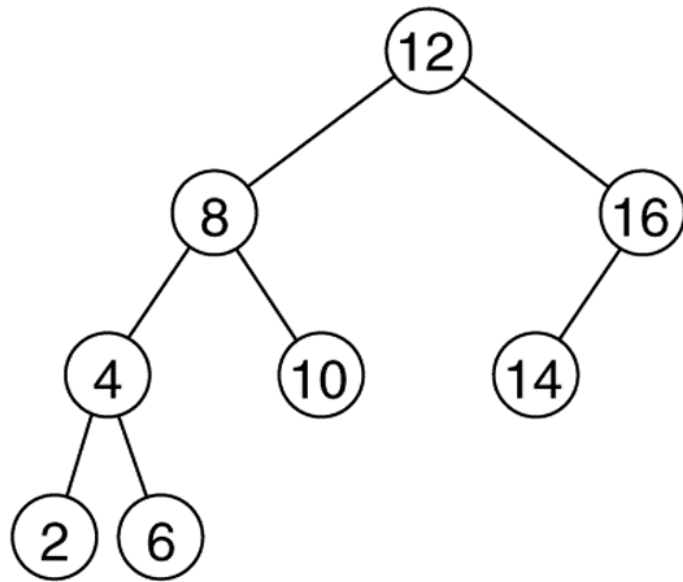


(e)

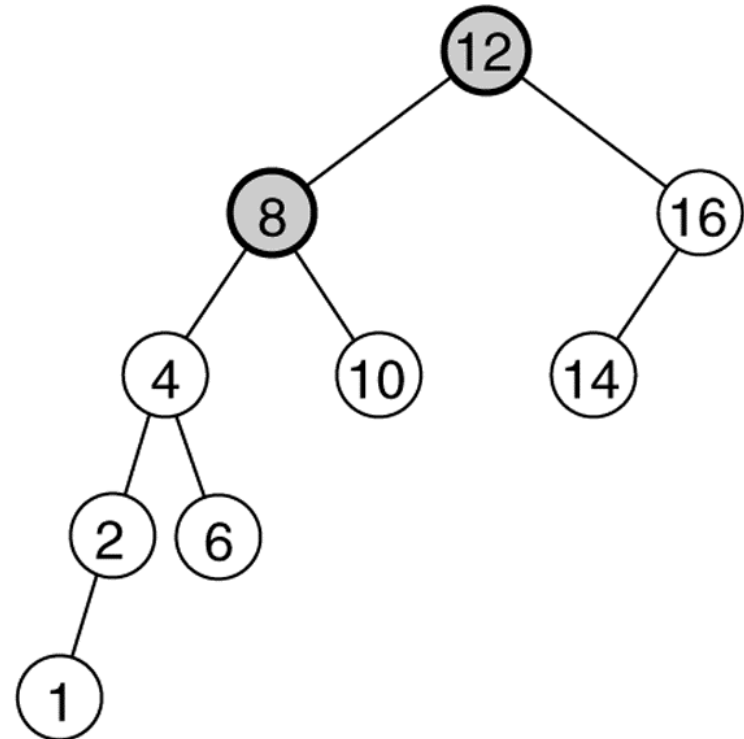


## Figure 19.21

Two binary search trees: (a) an AVL tree; (b) not an AVL tree (unbalanced nodes are darkened)



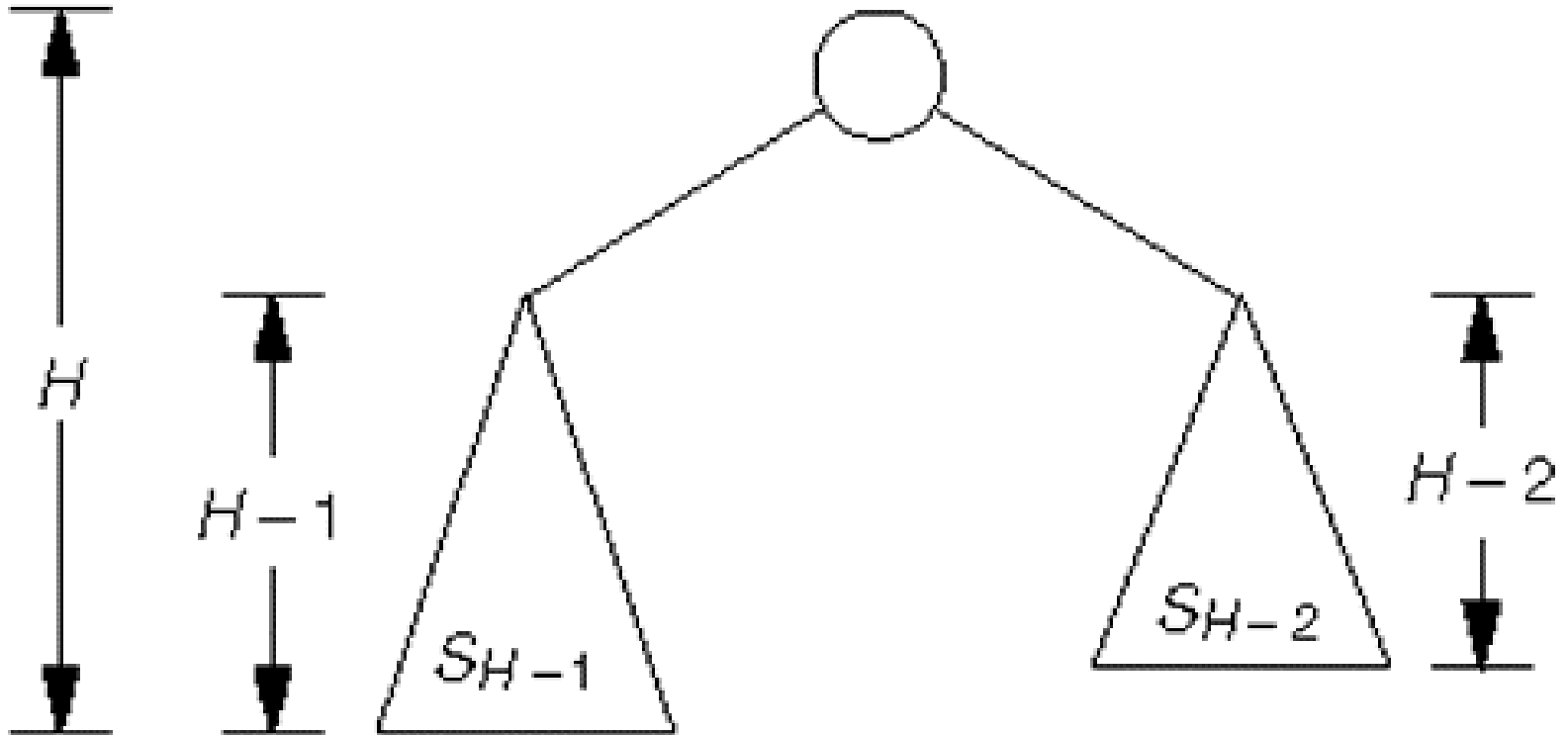
(a)



(b)

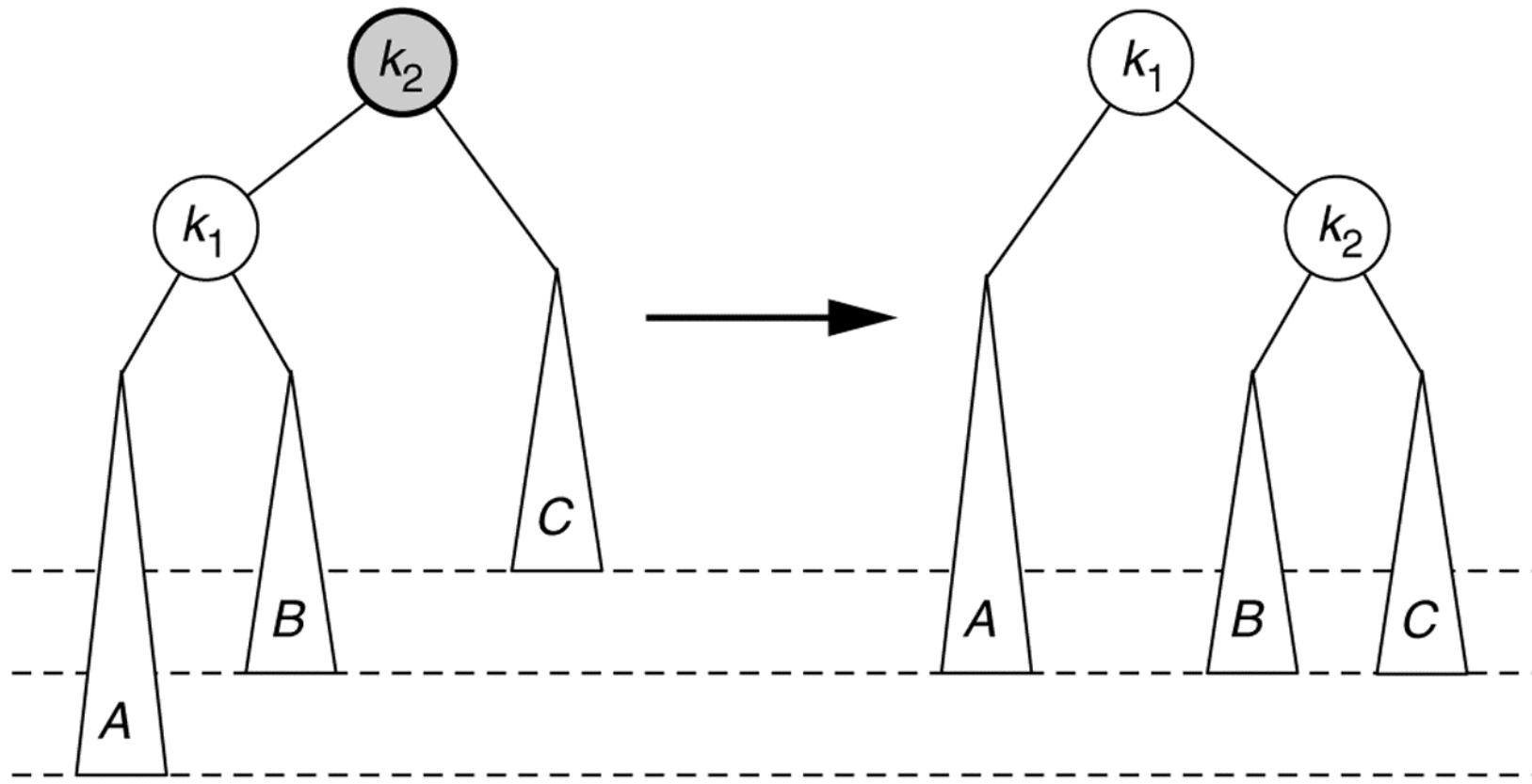
# Figure 19.22

Minimum tree of height  $H$



# Figure 19.23

Single rotation to fix case 1

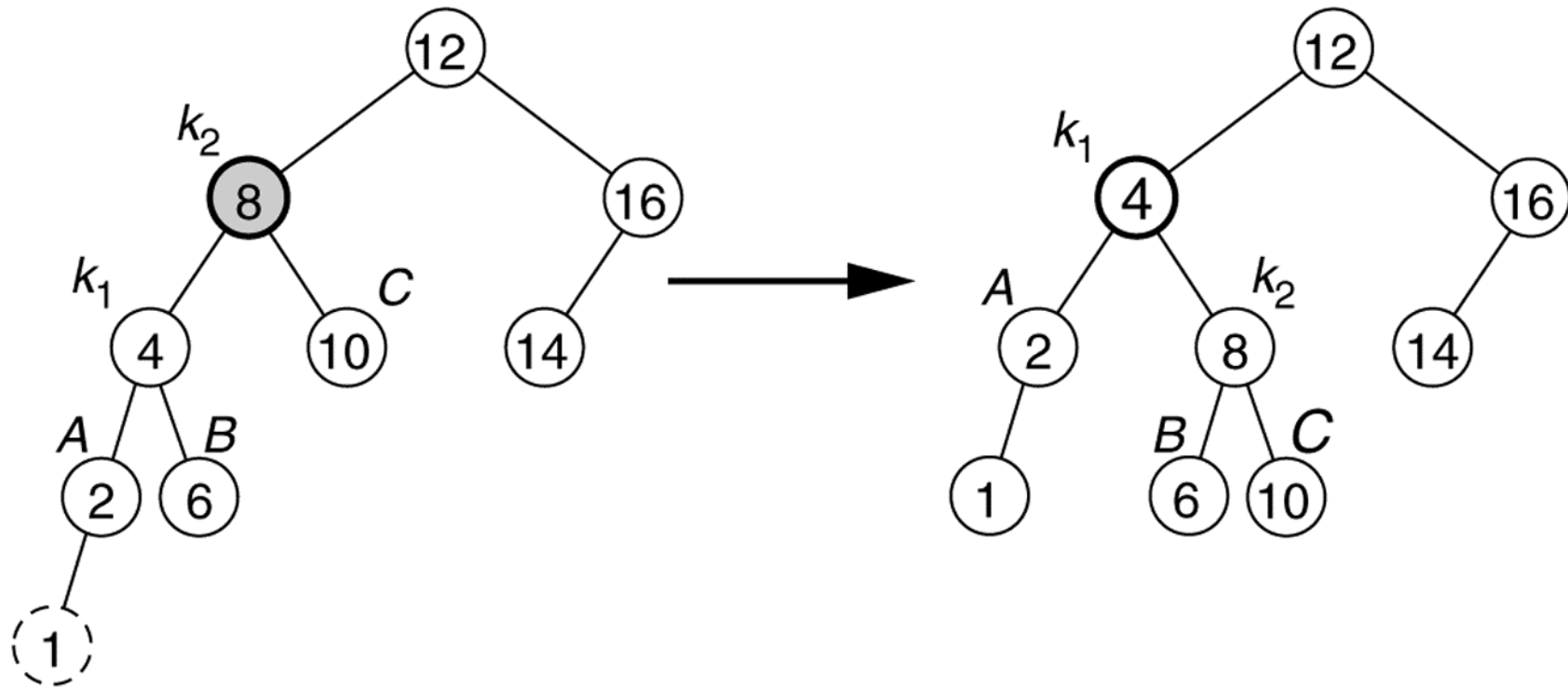


(a) Before rotation

(b) After rotation

# Figure 19.25

Single rotation fixes an AVL tree after insertion of 1.

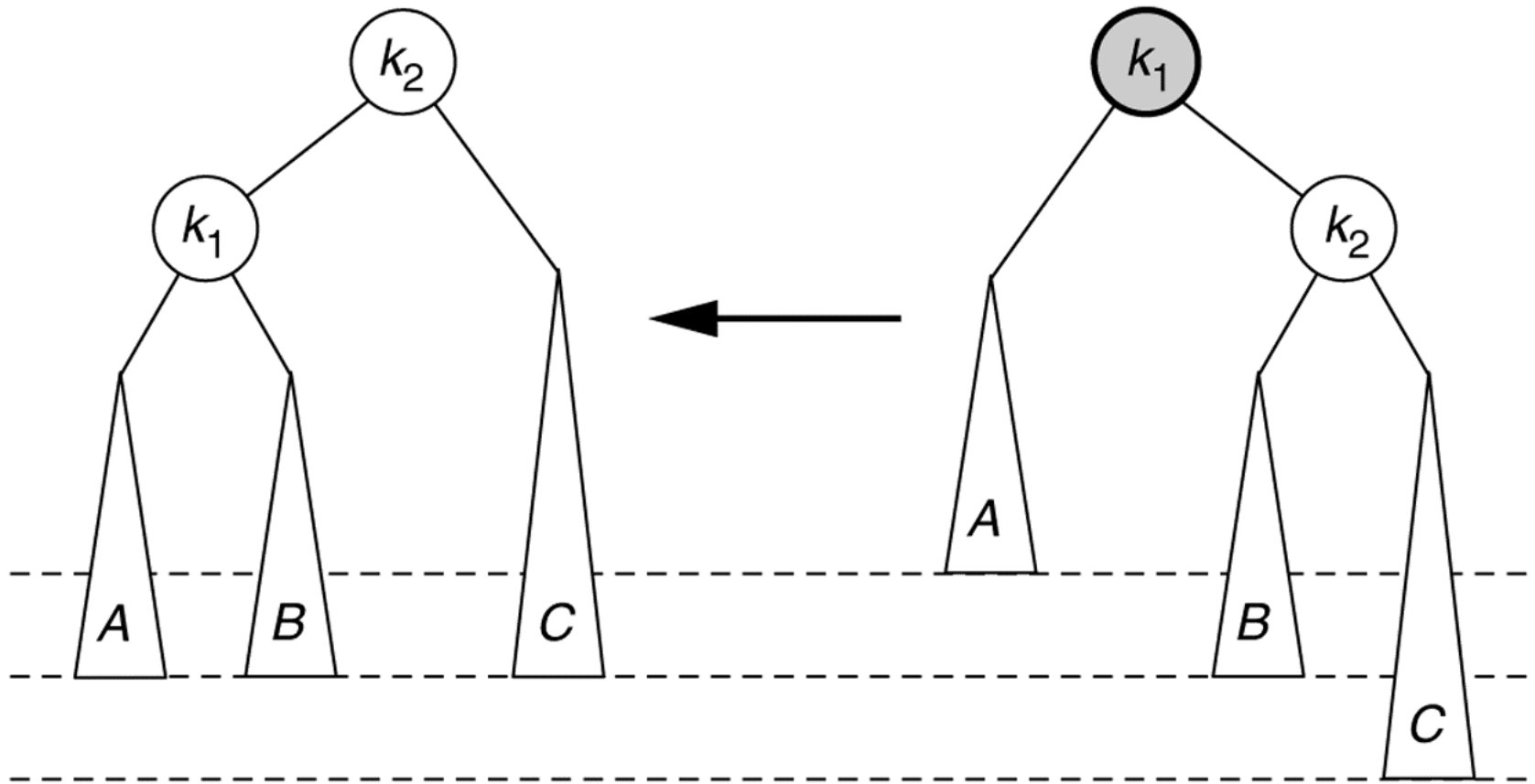


(a) Before rotation

(b) After rotation

# Figure 19.26

Symmetric single rotation to fix case 4

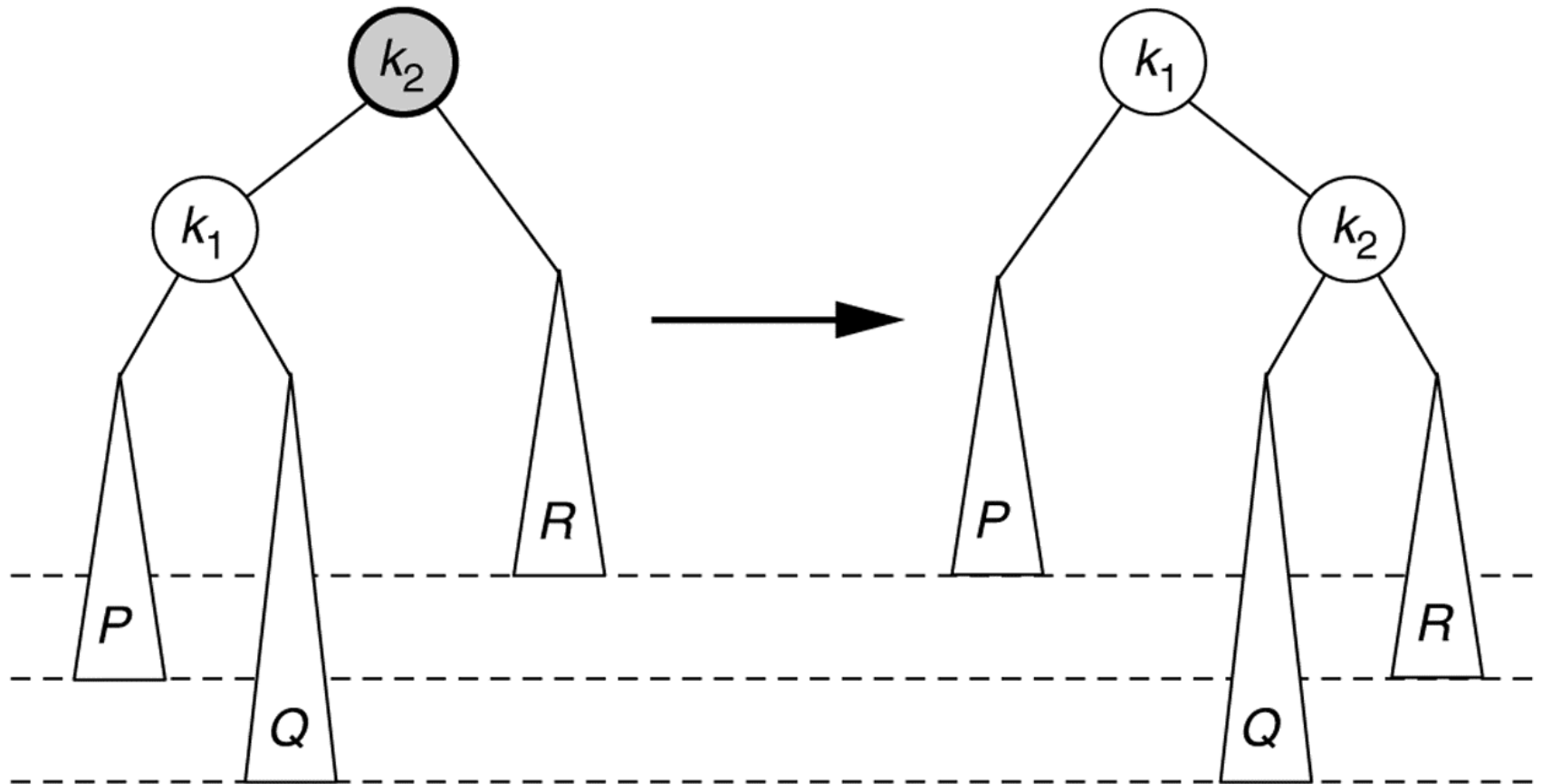


(a) After rotation

(b) Before rotation

# Figure 19.28

Single rotation does not fix case 2.

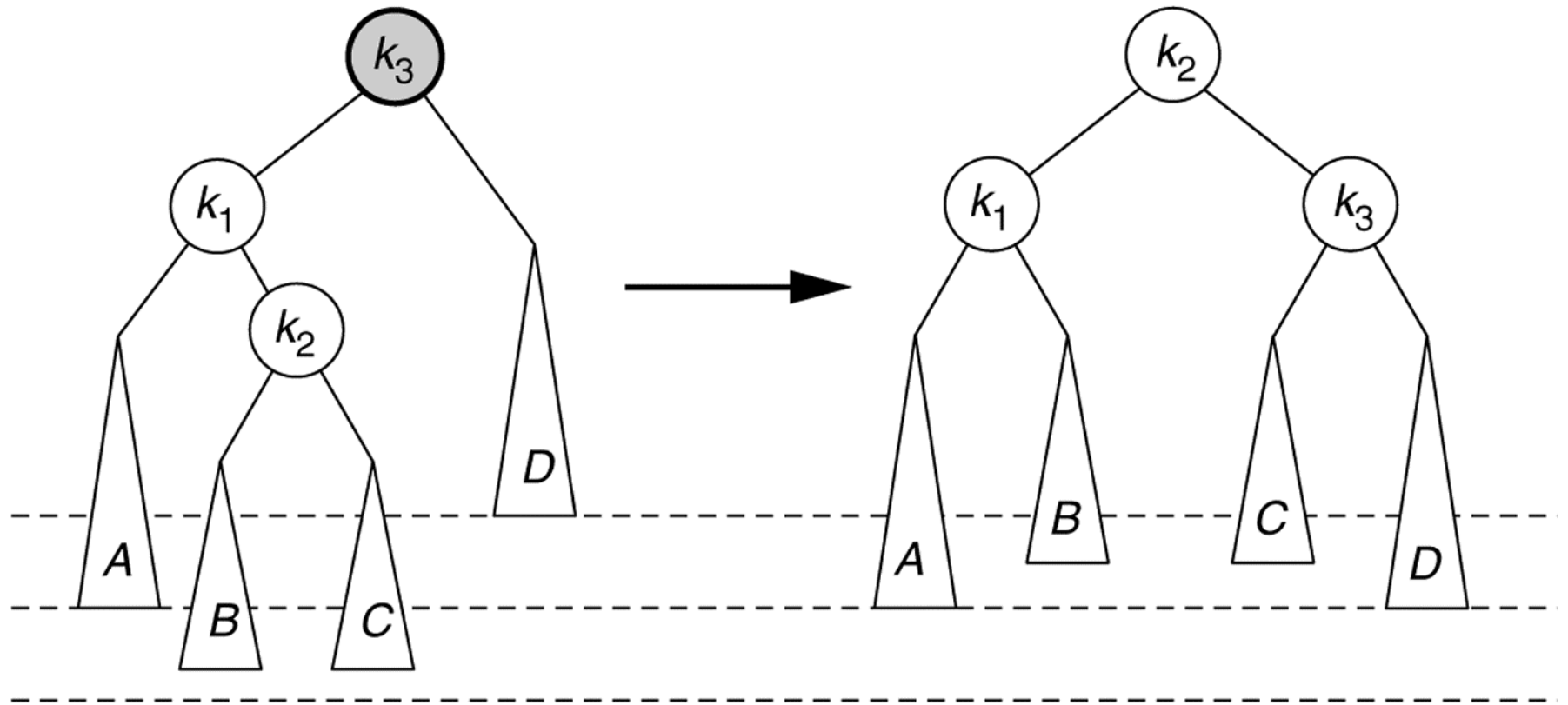


(a) Before rotation

(b) After rotation

# Figure 19.29

Left-right double rotation to fix case 2

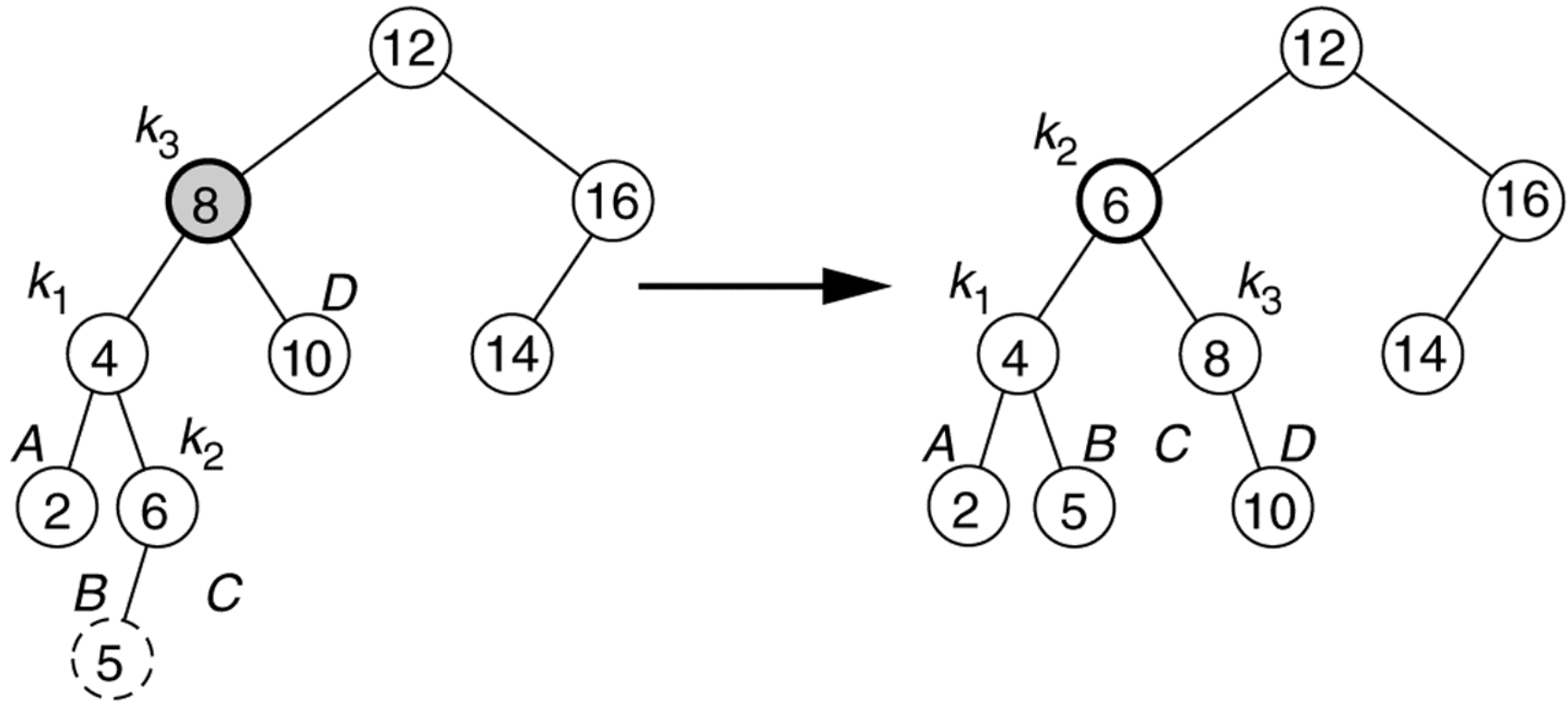


(a) Before rotation

(b) After rotation

# Figure 19.30

Double rotation fixes AVL tree after the insertion of 5.



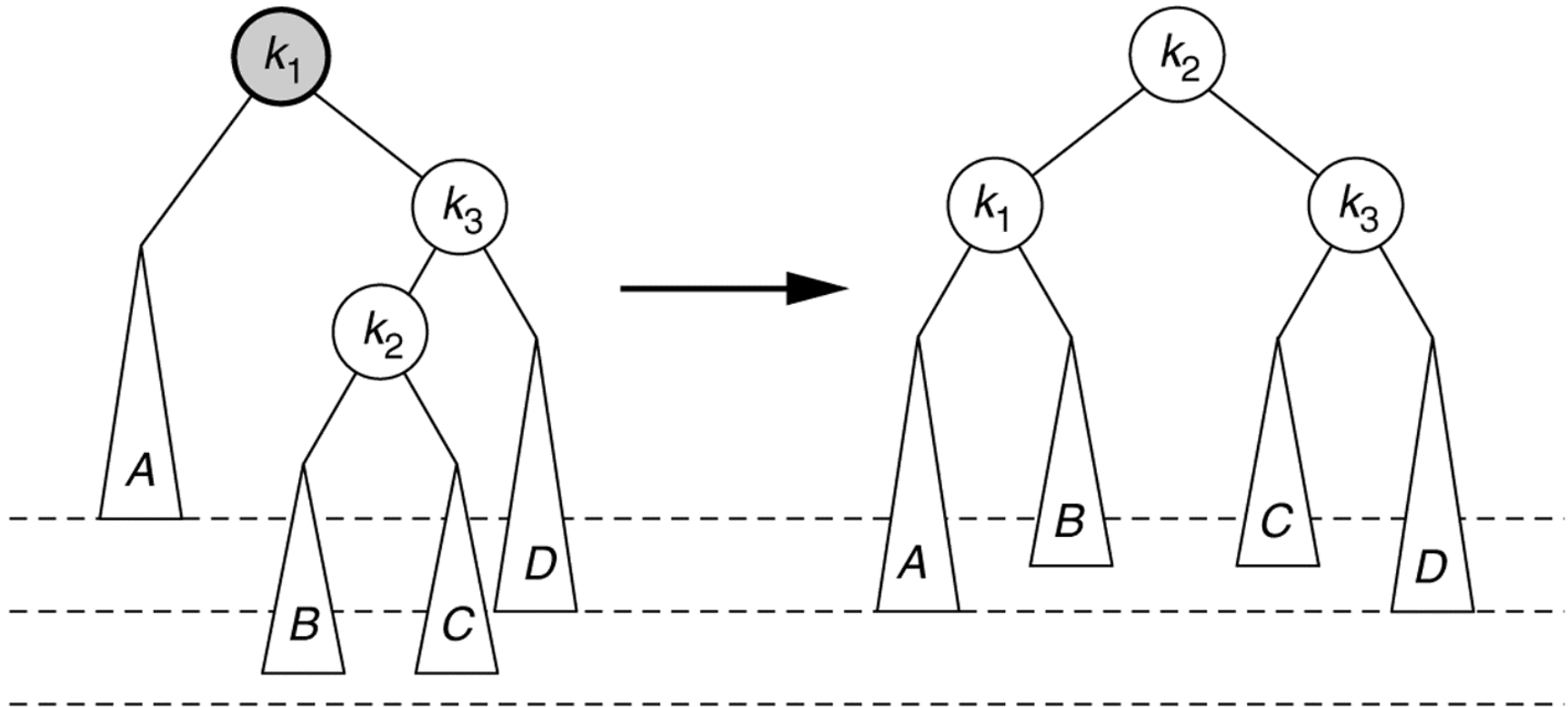
(a) Before rotation

(b) After rotation



# Figure 19.31

Right-Left double rotation to fix case 3.

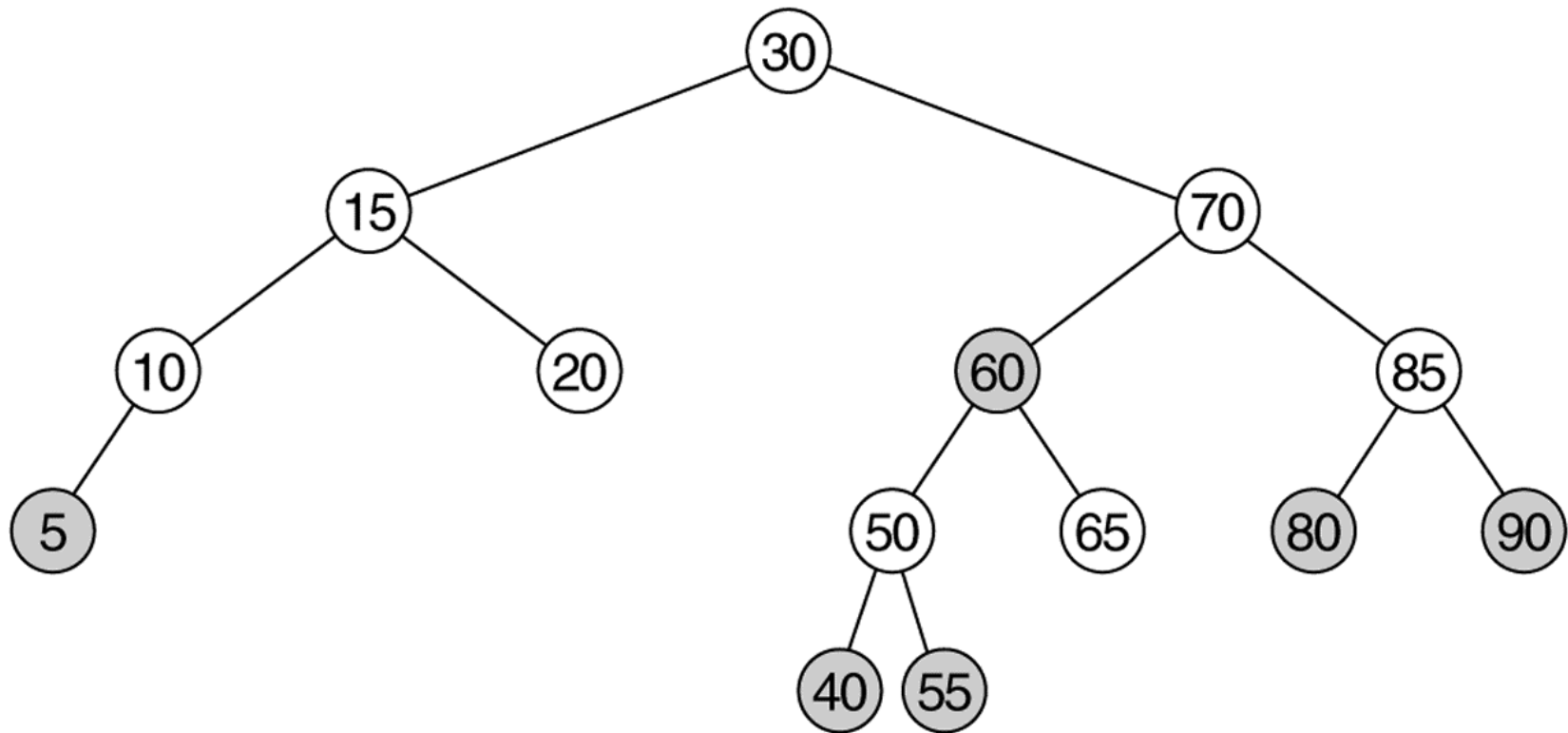


(a) Before rotation

(b) After rotation

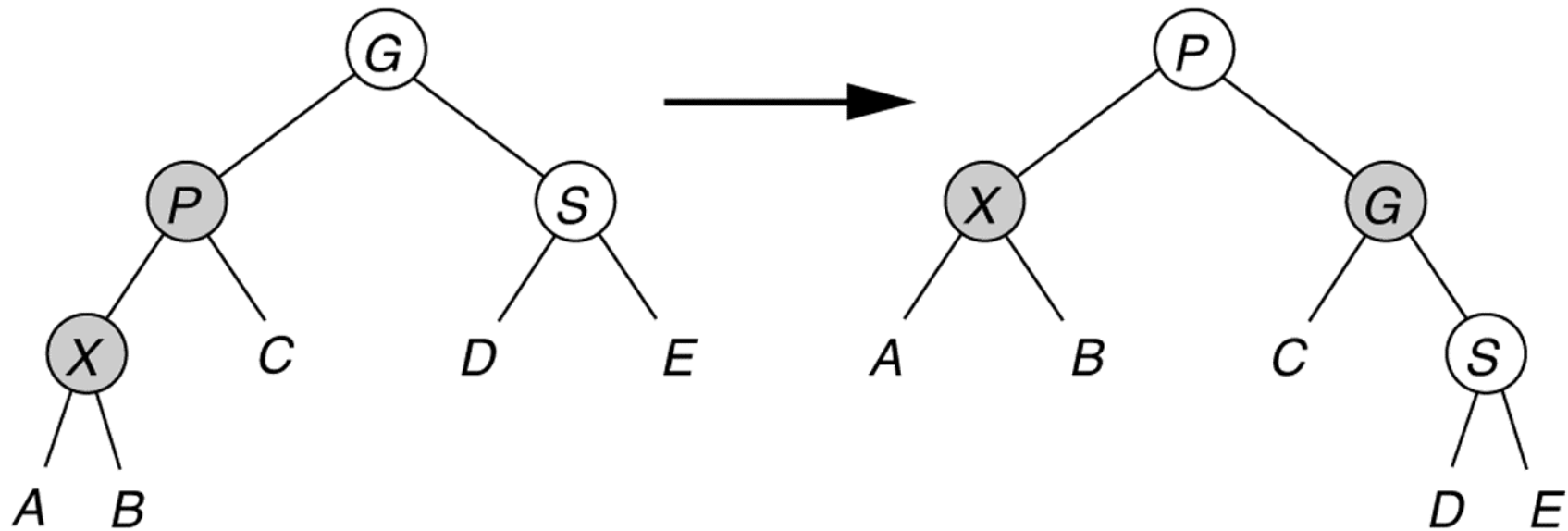
## Figure 19.34

A red–black tree: The insertion sequence is 10, 85, 15, 70, 20, 60, 30, 50, 65, 80, 90, 40, 5, and 55 (shaded nodes are red).



## Figure 19.35

If  $S$  is black, a single rotation between parent and grandparent, with appropriate color changes, restores property 3 if  $X$  is an outside grandchild.

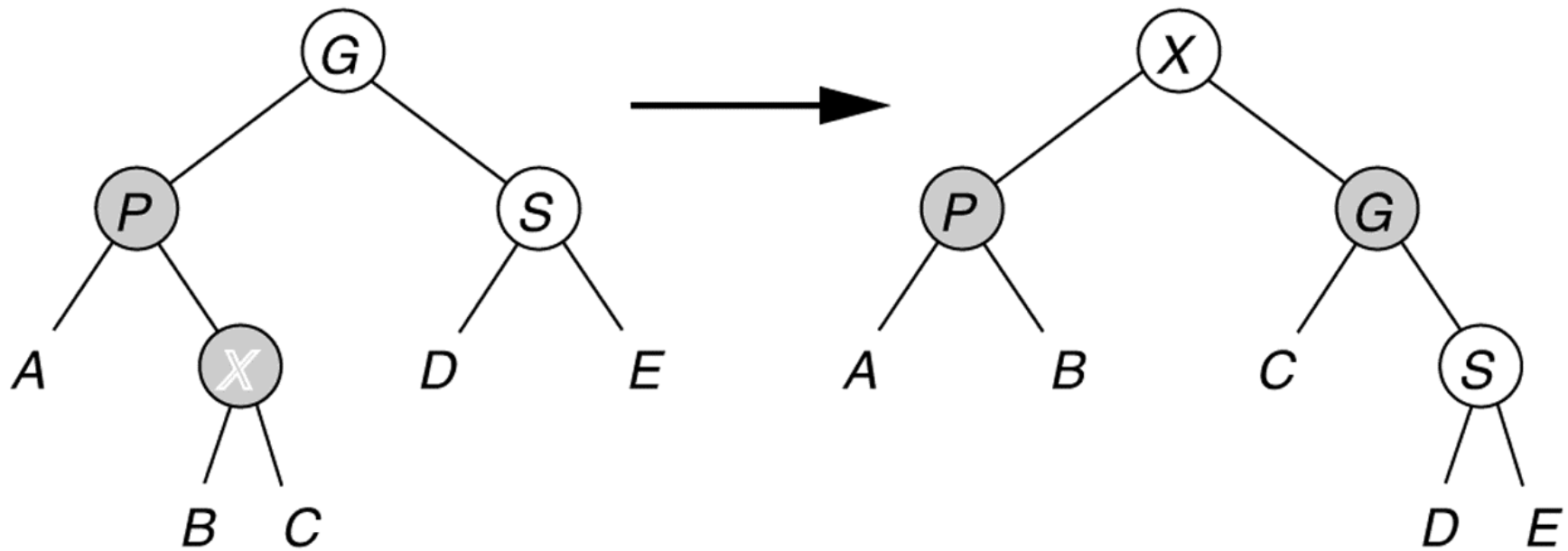


(a) Before rotation

(b) After rotation

## Figure 19.36

If  $S$  is black, a double rotation involving  $X$ , the parent, and the grandparent, with appropriate color changes, restores property 3 if  $X$  is an inside grandchild.

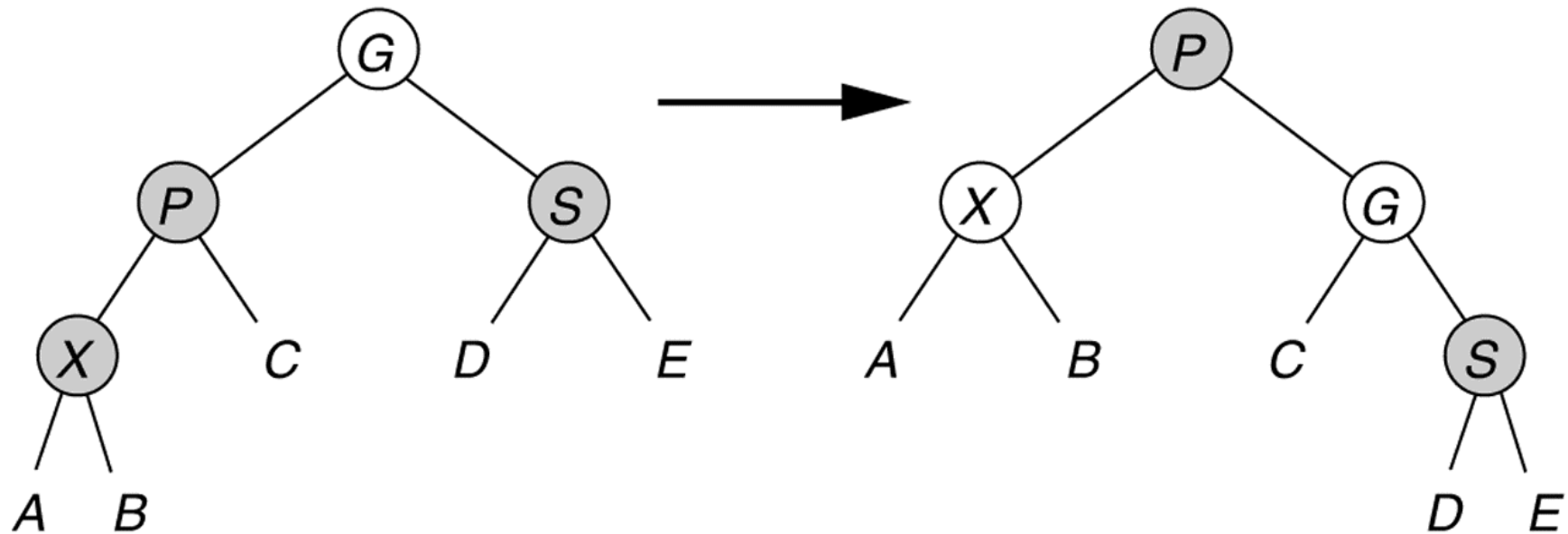


(a) Before rotation

(b) After rotation

## Figure 19.37

If  $S$  is red, a single rotation between parent and grandparent, with appropriate color changes, restores property 3 between  $X$  and  $P$ .



(a) Before rotation

(b) After rotation

## Figure 19.38

Color flip: Only if  $X$ 's parent is red do we continue with a rotation.

